

GPU Optimization of Advanced Encryption Standard Cihangir Tezcan, PhD Graduate School of Informatics, METU, Ankara



TÜBİTAK



Data Encryption Standard (DES)

- Designed by IBM in 1970s, based on an earlier design by Feistel.
- In 1976, NSA tweaked the algorithm by changing its S-boxes.
 - Block Size: 64 bits
 - Key Length: 56 bits
 - **Rounds:** 16
- Currently known as Data Encryption Algorithm (DEA) since it is no longer a standard.
- Became useless after 1990s since its short key is susceptible to brute force attacks.



Short Keys

- An attacker that captures a single ciphertext, can try to decrypt it with every possible key to check if it provides a meaningful plaintext.
- Such an attack is called exhaustive search or brute force attack.
- Exhaustive search is a generic attack, i.e. valid for every cipher.
- For a k-bit keyed cipher, the attacker is required to perform at most 2^k encryptions/decryptions.
- Thus, security of a block cipher is upper bounded by exhaustive search.



Short Keys

| 2 ⁵⁶ = | 72,057,594,037,927,936 |
|---------------------------|---|
| 2 ⁸⁰ = | 1,208,925,819,614,629,174,706,176 |
| 2 ¹²⁸ = | 340,282,366,920,938,463,463,374,607,431,768,211,456 |
| 2 ¹⁹² = | 6,277,101,735,386,680,763,835,789,423,207,666,416,102,355,444,464,034,512,896 |
| 2 ²⁵⁶ = | 115,792,089,237,316,195,423,570,985,008,687,907,853,269,984,665,640,564,039,457,584,007,913,129,639,936 |



Advanced Encryption Standard (AES)

- Rijndael is designed by Joan Daemen and Vincent Rijmen.
- Standardized in 2001 by NIST (winner of the AES competition) and named AES.
 - Block Size: 128 bits
 - Key Length: 128, 192, 256 bits
 - Rounds: 10, 12, 14 (depends on the key length)
 - Type: SPN
- Known attacks are ineffective.



Byte ordering of AES

| 0 | 4 | 8 | 12 | | |
|---|---|----|----|--|--|
| 1 | 5 | 9 | 13 | | |
| 2 | 6 | 10 | 14 | | |
| 3 | 7 | 11 | 15 | | |

A byte can be treated in two ways when dealing with AES:

- 1. A byte can be viewed as a string of 8 bits
- 2. A byte can be viewed as an element of Galois Field GF(2⁸)



Arithmetic in *GF(2⁸)*

- AES uses the Galois Field defined by the irreducible polynomial $R_p = X^8 + X^4 + X^3 + X + 1$
- A byte can be represented by a degree 7 polynomial where the bits of the byte corresponds to coefficients of this polynomial.
- Now we can perform addition and multiplication on these polynomials modulo $R_{\rm p}$
- This more mathematical representation is better to understand the security of the cipher



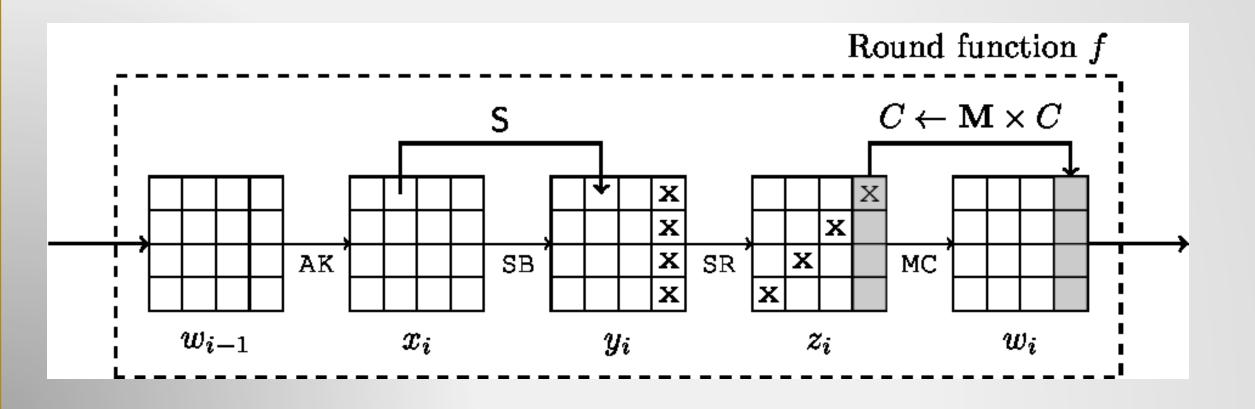
Round Function of AES

Round function of AES contains the following operations:

- 1. Key Addition (AK): XOR the 128-bit round key with the input
- 2. Sub Bytes (SB): Apply 8x8 S-box on 16 bytes
- 3. Shift Rows (SR): Rotate rows to the left
- 4. Mix Columns (MC): Multiply by matrix M



Advanced Encryption Standard (AES)





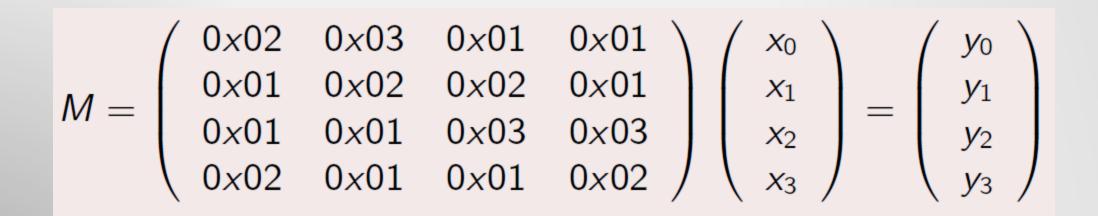
| | 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9 | а | b | С | d | е | f |
|---|----|----|-----------------|----|----|----|----|----|------------------|----|----|----|----|------------|----|----|
| 0 | 63 | 7c | 77 | 7b | f2 | 6b | 6f | c5 | 30 | 01 | 67 | 2b | fe | d7 | ab | 76 |
| 1 | са | 82 | c9 | 7d | fa | 59 | 47 | fO | ad | d4 | a2 | af | 9c | a4 | 72 | c0 |
| 2 | b7 | fd | 93 | 26 | 36 | 3f | f7 | сс | 34 | a5 | e5 | f1 | 71 | d 8 | 31 | 15 |
| 3 | 04 | c7 | 23 | c3 | 18 | 96 | 05 | 9a | 07 | 12 | 80 | e2 | eb | 27 | b2 | 75 |
| 4 | 09 | 83 | 2c | 1a | 1b | 6e | 5a | a0 | 52 | 3b | d6 | b3 | 29 | e3 | 2f | 84 |
| 5 | 53 | d1 | 00 | ed | 20 | fc | b1 | 5b | 6a | cb | be | 39 | 4a | 4c | 58 | cf |
| 6 | d0 | ef | aa | fb | 43 | 4d | 33 | 85 | 45 | f9 | 02 | 7f | 50 | 3c | 9f | a8 |
| 7 | 51 | a3 | 40 | 8f | 92 | 9d | 38 | f5 | bc | b6 | da | 21 | 10 | ff | f3 | d2 |
| 8 | cd | 0c | 13 | ec | 5f | 97 | 44 | 17 | c4 | a7 | 7e | 3d | 64 | 5d | 19 | 73 |
| 9 | 60 | 81 | 4f | dc | 22 | 2a | 90 | 88 | 46 | ee | b8 | 14 | de | 5e | 0b | db |
| а | e0 | 32 | 3a | 0a | 49 | 06 | 24 | 5c | c2 | d3 | ac | 62 | 91 | 95 | e4 | 79 |
| b | e7 | c8 | 37 | 6d | 8d | d5 | 4e | a9 | <mark>6</mark> c | 56 | f4 | ea | 65 | 7a | ae | 08 |
| с | ba | 78 | 25 | 2e | 1c | a6 | b4 | c6 | e8 | dd | 74 | 1f | 4b | bd | 8b | 8a |
| d | 70 | 3e | b5 | 66 | 48 | 03 | f6 | 0e | 61 | 35 | 57 | b9 | 86 | c1 | 1d | 9e |
| е | e1 | f8 | <mark>98</mark> | 11 | 69 | d9 | 8e | 94 | 9b | 1e | 87 | e9 | ce | 55 | 28 | df |
| f | 8c | a1 | 89 | 0d | bf | e6 | 42 | 68 | 41 | 99 | 2d | Of | b0 | 54 | bb | 16 |

S-box of AES

Input 3b means the 3rd row and bth column (i.e. $S(3b) = e^2$)



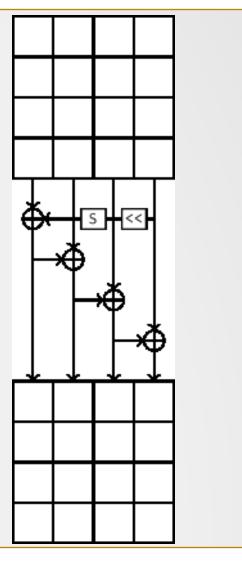
Mix Columns (MC) (Omitted in the last round!)





Key Schedule

- The 128-bit master key is also the first round key
- Next round key is obtained by the previous one by performing:
 - 1. Byte rotation
 - 2. Four S-box operation
 - 3. Round Constant Addition
 - 4. XOR of 32-bit values





Mode of Operations for Block Ciphers

- When we use a block cipher, we need to choose a mode of operation to determine how we encrypt plaintext that is larger than a single block.
- Some mode of operations are not secure.
- Some mode of operations allow parallelization.





Thanks



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